

# Some ideas for a finite geometry of interaction model of second-order MLL

NGUYỄN Lê Thành Dũng (LIPN, Université Paris 13)

nlttd@nguyentito.eu

In a recent paper<sup>1</sup>, we showed that in second-order Multiplicative Linear Logic (MLL2), the cut-elimination between a proof net  $\pi : A$  and a “test”  $\rho : A \multimap B$  ( $B$  quantifier-free) can be performed as a sort of “dialogue”: the two proofs exchange positions of axiom links whose endpoints are atoms of  $A$ . The observational equivalence class of  $\pi$  can therefore be determined by a “strategy” of bounded size, proving that the observational quotient of MLL2 is finite at each type. However, the ad-hoc notion of strategy used does not compose.

An attempt to build an actual denotational semantics based on these ideas is currently underway. A first idea is to use as “strategies” *stable functions*  $L_{\exists} \rightarrow L_{\forall}$  where  $L_{\exists}$  (resp.  $L_{\forall}$ ) is a coherence space whose cliques correspond to (partial) axiom linkings over the existential (resp. universal) atoms of  $A$ . Then, from the strategies of  $\pi$  and  $\rho$ , the quantifier-free observation can be characterized by a “feedback equation” à la Geometry of Interaction (GoI).

Next, we consider the compact closed category  $\mathcal{C} = \text{Int}(\text{CohS})$  i.e. the Joyal–Street–Verity  $\text{Int}$  construction applied to the category  $\text{CohS}$  of coherence spaces and stable functions. It is a “wave-style” GoI model of propositional MLL, and the goal is to interpret quantifiers so that the semantics of closed types fit with the specifications of the previous paragraph. This would give a finite semantics for MLL2.

It seems that this model (which, at the time of writing, still depends on unproven lemmas!) exhibits a phenomenon of *inversion of polarities*: the universal quantifier behaves like a positive connective. Indeed, a universal proof plays axiom links, while an existential proof reacts to them (to compare with Hughes’s hypergames, it is as if, after a copycat link is first set up by a polymorphic program, memorizing it for the rest of the interaction were the responsibility of the user of this program). This is reminiscent of the impredicative encodings of data types<sup>2</sup>, with initial algebras having universal types and final coalgebras having existential types.

---

<sup>1</sup>*Finite semantics of polymorphism, complexity and the power of type fixpoints*, with Paolo Pistone, Thomas Seiller and Lorenzo Tortora de Falco. Available at <https://hal.archives-ouvertes.fr/hal-01979009/document>.

<sup>2</sup>To quote J.-Y. Girard (*The Blind Spot*, §12.B.2): “The second-order translations do invert polarities. [...] This inversion of polarities is due to deep reasons – so deep that they remain completely obscure.”